

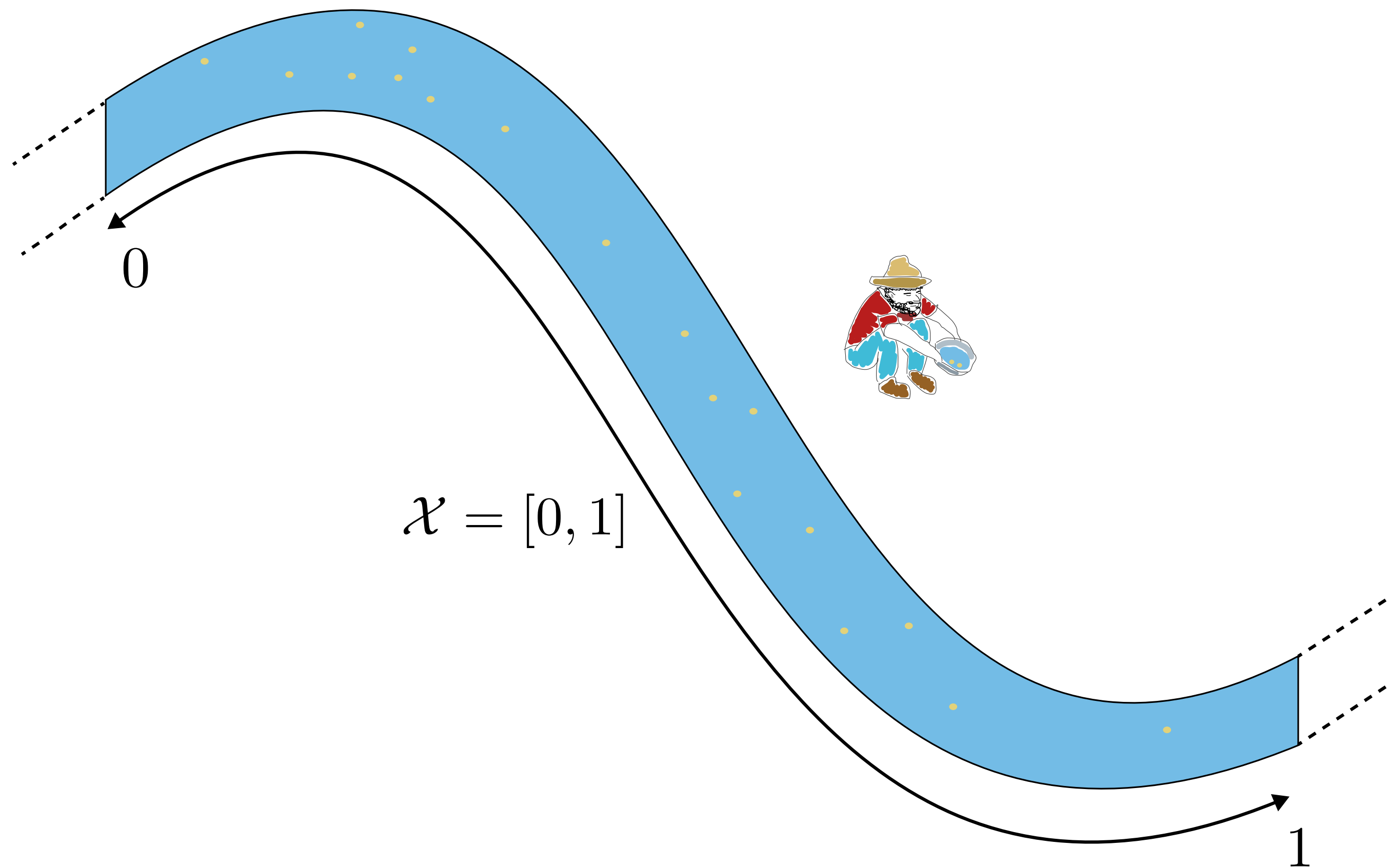
# On some adaptivity questions in multi-armed bandits

Thèse de doctorat (virtually) at Orsay

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# Gold panning in a river



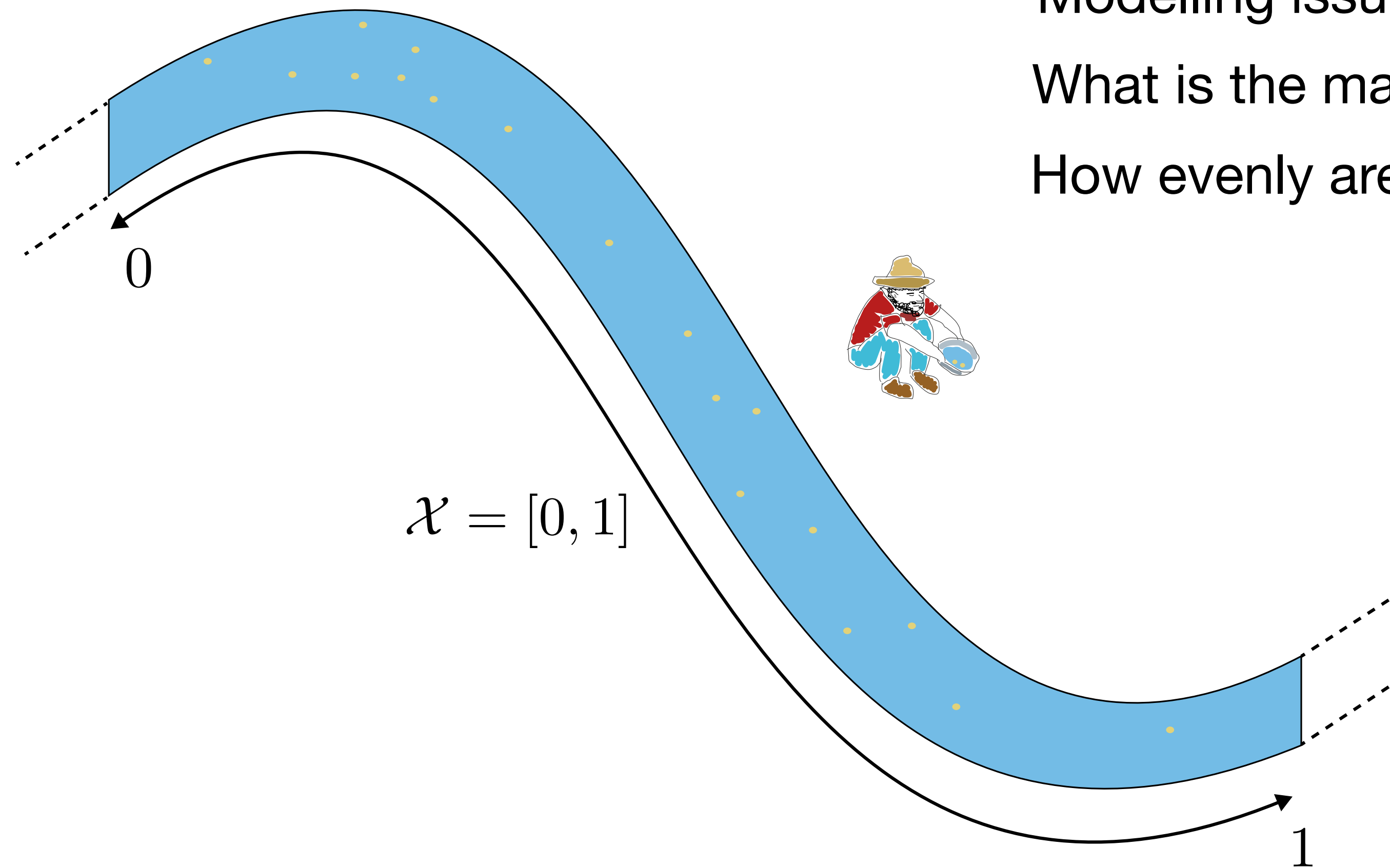
Every day  $t = 1, \dots,$

- Pick a spot  $X_t \in [0, 1]$  on the river
- Get  $Y_t$  grams of gold

What strategy to  
get as much gold as possible?

Making as few assumptions  
as necessary

# Gold panning in a river



Modelling issues:

What is the maximum reward one can get in a single round?

How evenly are the rewards distributed across space?

Should we just guess?

This will crucially affect our decisions!

what is the maximum reward I can get in a single round?



how evenly are the rewards distributed across space?



**How does prior information about the rewards  
affect optimal strategies?**

# Outline

- I/ Multi-armed bandits
- II/ Adapting to the unknown range of the rewards
- III/ Continuous bandits and smoothness

# Multi-armed bandits

[Thompson 1933] [Robbins 1952]

$K$  probability distributions  $(\nu_1, \dots, \nu_K)$  unknown to the player

$$\mu_i = \mathbb{E}(\nu_i)$$

At every time-step  $t$  the player:

- chooses an action  $A_t$
- receives and observes reward  $Y_t \sim \nu_{A_t}$  given  $A_t$

Minimize **regret**

$$R_T = T\mu^* - \mathbb{E}\left[\sum_{t=1}^T Y_t\right] = T\mu^* - \mathbb{E}\left[\sum_{t=1}^T \mu_{A_t}\right]$$

Prior knowledge = assumption on  $\nu_1, \dots, \nu_K$

# Upper-Confidence Bounds

[Auer, Cesa-Bianchi, Fischer, '02]

$\nu_1, \dots, \nu_K$  are supported in  $[0, 1] \Leftrightarrow$  all rewards are in  $[0, 1]$

Draw every arm once

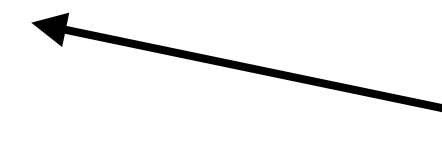
Then draw arm  $A_t$  maximizing the index

$$U_a(t) = \hat{\mu}_a(t) + \sqrt{\frac{2 \log t}{N_a(t)}}$$

Empirical mean of rewards from arm  $a$



Number of times  $a$  was picked



## Distribution-free

$$R_T \leq c\sqrt{KT \ln T} + K$$

## Distribution-dependent

$$R_T \leq \sum_{a=1}^K \frac{8}{\mu^* - \mu_a} \log T + \sum_{a=1}^K (\mu^* - \mu_a)$$

# Improving UCB: two types of optimality

UCB explores a bit too much: we can reduce the confidence bounds and get better guarantees

**MOSS** [Audibert, Bubeck '09]  
[Degenne, Perchet '16]

$$R_T \leq c\sqrt{KT \cancel{\ln T}} + K$$

**KL-UCB** [Cappé, Garivier, Maillard, Munos, Stoltz '13]

$$R_T \leq \sum_{a=1}^K \frac{\mu^* - \mu_a}{\mathcal{K}_{\text{inf}}(\nu_a, \mu^*)} \log T + \text{cst}$$

$$\mathcal{K}_{\text{inf}}(\nu_a, \mu^*) = \inf \{ \text{KL}(\nu_a, \nu') \mid \mathbb{E}(\nu') > \mu^* \text{ and } \nu'([0, 1]) = 1 \}$$

None are further improvable [Auer, Cesa-Bianchi, Freund, Schapire '02], [Lai & Robbins 1985]

[Garivier, Ménard '17] obtain joint optimality in a parametric setting

**Theorem: KL-UCB-switch [Garivier, H, Ménard, Stoltz '18]**

KL-UCB-switch is both distribution-dependent and distribution-free optimal

# What if the range is unknown?

rewards are in  $[m, M]$  instead of  $[0, 1]$ , and  $m$  and  $M$  are unknown to the player?

# Unknown range: initial remarks

1.  $\frac{Y_t - m}{M - m} \in [0, 1]$

playing with rescaled rewards is equivalent  
to standard game with regret scaled by  $(M-m)$   
... but we cannot rescale

2. (KL-)UCB, MOSS, etc. are scale-dependent

$$\hat{\mu}_a(t) + (M - m) \sqrt{\frac{\log t}{N_a(t)}}$$

3. Not knowing the range is harder than knowing the range

$$\sup_{\text{Problems in } [m, M]} R_T \geq c (M - m) \sqrt{KT}$$

Can we match this lower bound without knowing  $m$  and  $M$ ?

# Unknown range: distribution-free adaptation

AdaHedge is from de [Rooij, van Erven, Grünwald, Koolen '13]

[Cesa-Bianchi, Mansour, Stoltz '07]

## Theorem: Minimax range adaptation (H, Stoltz 2020)

AdaHedge for bandits with extra-exploration, guarantees

$$\text{for all } m < M, \quad \sup_{\text{Prob in } [m, M]} R_T \leq 7(M-m)\sqrt{TK \ln K} + 10(M-m)K \ln K$$

Same guarantees as if we had known the range in advance!

But can we get distribution-dependent  $\log T$  bounds?

Obstacle: extra-exploration forces at least  $R_T \geq (M-m)\sqrt{KT}$

# Adaptive rates

## Distribution-free rate

For all  $T$ , for all  $M > 0$ ,  $\sup_{\text{Problems in } [0, M]} R_T \leq M \Phi_{\text{free}}(T)$

e.g.

Bandit AdaHedge

enjoys

$$\Phi_{\text{free}}(T) = \sqrt{TK \ln K}$$

## Distribution-dependent rate

For all  $M > 0$ , for all problems in  $[0, M]$ ,  $\limsup_{T \rightarrow \infty} \frac{R_T}{\Phi_{\text{dep}}(T)} < +\infty$

e.g.

$$\tilde{U}_a(t) = \hat{\mu}_a(t) + (\log t) \sqrt{\frac{\log t}{N_a(t)}}$$

enjoys

$$\Phi_{\text{dep}}(T) = (\log T)^2 \quad [\text{Lattimore '17}]$$

Can we get  $\log T$  and  $\sqrt{KT}$  simultaneously?

# Lower bound for adaptation to the range

**Theorem: Lower bound for range adaptation (H, Stoltz 2020)**

For any algorithm enjoying rates  $\Phi_{\text{free}}(T)$  and  $\Phi_{\text{dep}}(T)$

$$\Phi_{\text{dep}}(T) \Phi_{\text{free}}(T) \gtrsim T$$

in particular,  $\Phi_{\text{free}}(T) \leq \mathcal{O}(\sqrt{T}) \Rightarrow \Phi_{\text{dep}}(T) \geq \Omega(\sqrt{T})$

The cost of distribution-free adaptation is hidden in distribution-dependent rates!

We get  $T^\alpha$  and  $T^{1-\alpha}$  (for  $1/2 < \alpha < 1$ ) by tuning the extra-exploration in Bandit AdaHedge

# Adapting to the smoothness

# Back to the river: continuous bandits

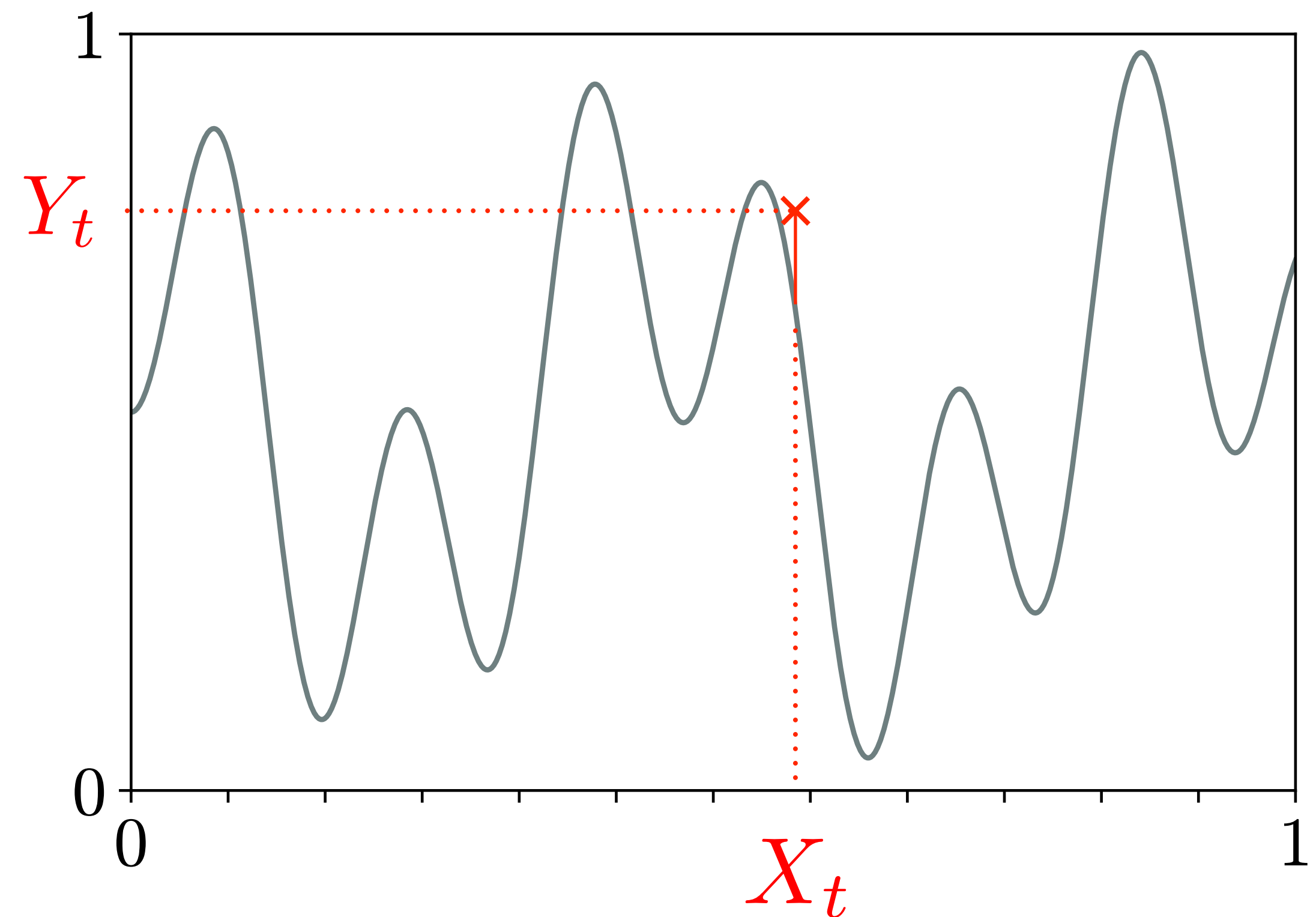
$\mathcal{X} = [0, 1]$     Unknown payoff distributions  $(\nu_x, x \in \mathcal{X})$     Mean-payoff function  $f : x \mapsto \mathbb{E}(\nu_x)$

For  $t = 1, \dots, T, \dots$ :

- pick  $X_t \in \mathcal{X}$
- receive and observe  $Y_t \sim \nu_{X_t}$  given  $X_t$

**Goal**

$$R_T = T \max_{x \in \mathcal{X}} f(x) - \mathbb{E} \left[ \sum_{t=1}^T f(X_t) \right]$$



We need assumptions

# Assumptions

1. Rewards are bounded in  $[0, 1]$

2. Mean-payoff function is “ $(L, \alpha)$ -Hölder around its max”

$$f(x^*) - f(x) \leq L |x^* - x|^\alpha$$

Lots of similar and/or more refined assumptions in the literature

[Agrawal 1995], [Kleinberg '05], [Auer, Ortner, Szepesvári, '07], [Bubeck, Munos, Stoltz, Szepesvári '11]  
[Bubeck, Stoltz, Yu '11], [Kleinberg, Slivkins, Upfal, '13], [Bull '15]

Related problems: bandit optimization/simple regret, maximum estimation, maximum location

$$\begin{array}{l} \text{find } \hat{X} \\ \text{s.t. } f(\hat{X}) \approx \max f \end{array}$$

[Valko, Carpentier, Munos '13]

[Bartlett, Gabillon, Valko '19] [Shang, Kauffman, Valko '19]

$$\begin{array}{c} \uparrow \\ \text{estimate } \max f \end{array}$$

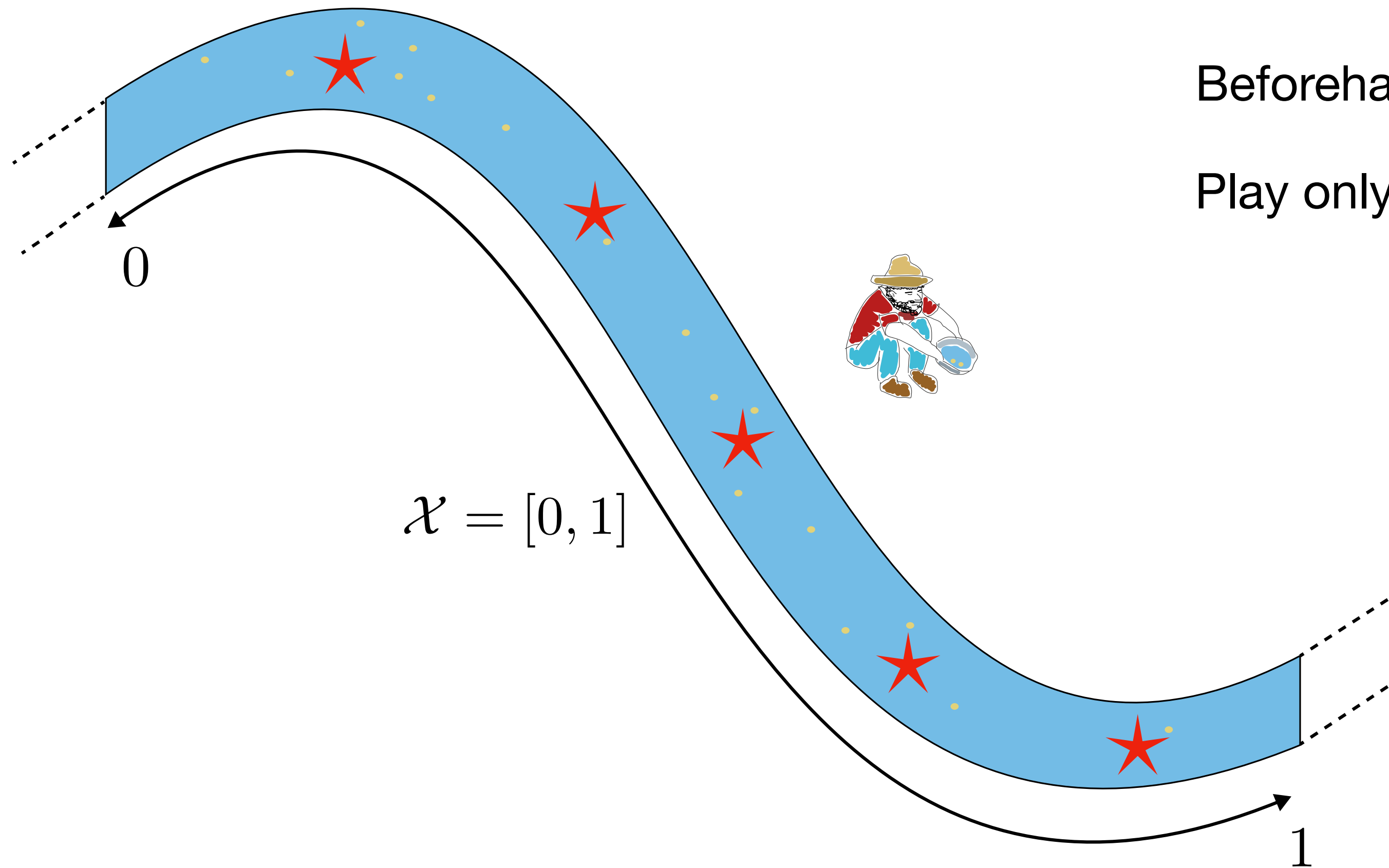
[Lepski 1994]

$$\begin{array}{c} \nwarrow \\ \text{find } \hat{X} \\ \text{s.t. } \hat{X} \approx x^* \end{array}$$

[Muller 1989]

# Discretization

[Kleinberg '05]

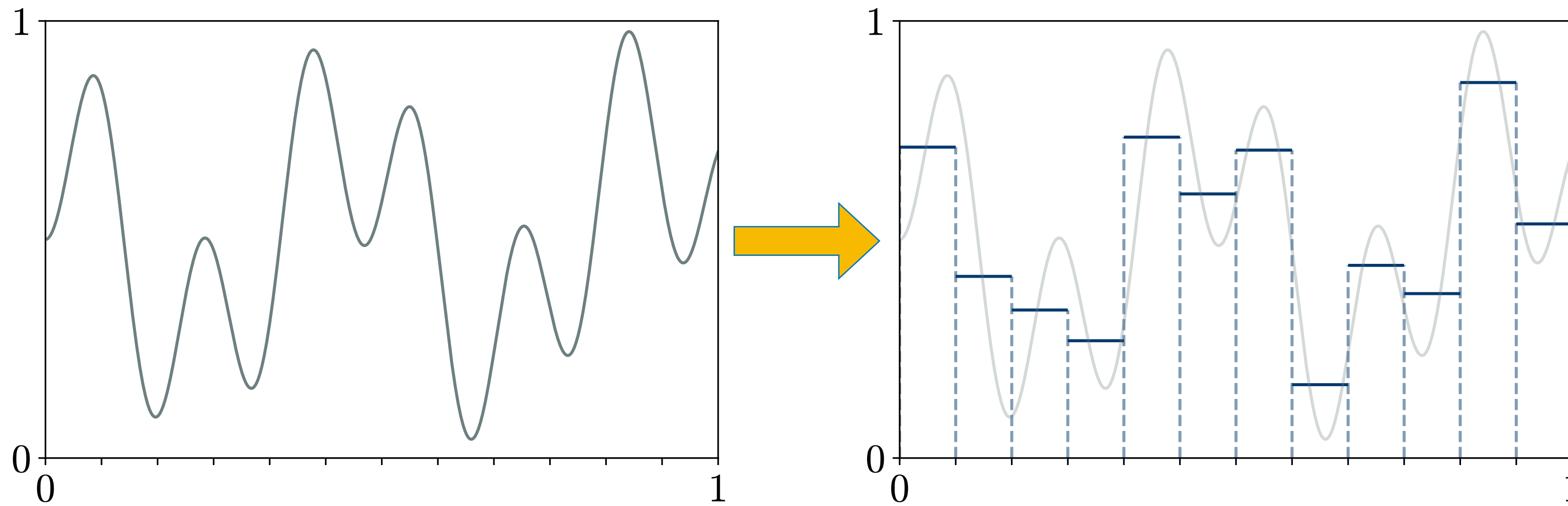


Beforehand, pick a finite number of spots ★

Play only at these spots, using a K-armed bandit algorithm

Back to a finite-armed bandit problem

# Bounding the regret of discretization



$$R_T \leqslant c L^{1/(2\alpha+1)} T^{(\alpha+1)/(2\alpha+1)} \quad \text{by tuning } K = K^*(\alpha) \text{ appropriately}$$

*Proof:*

$$R_T = T \left( \max f - \max_{1 \leqslant i \leqslant K} f(x_i) \right) + \max_{1 \leqslant i \leqslant K} f(x_i) - \mathbb{E} \left[ \sum_{t=1}^T f(X_t) \right]$$

$$\leqslant T \frac{L}{K^\alpha} + c\sqrt{KT}$$

# Impossibility of (full) adaptation

## Adaptive rates

An algorithm achieves adaptive rates  $\theta : \mathbb{R}_+ \rightarrow [1/2, 1]$  if

$$\sup_{f \text{ } \alpha\text{-H\"older}} R_T \leq c T^{\theta(\alpha)} \text{ for all } \alpha$$

Question: how can we obtain adaptive rates  $\theta(\alpha) = \frac{\alpha + 1}{2\alpha + 1}$ ?      Model selection? Cross-validation?  
... **Exploration is costly**

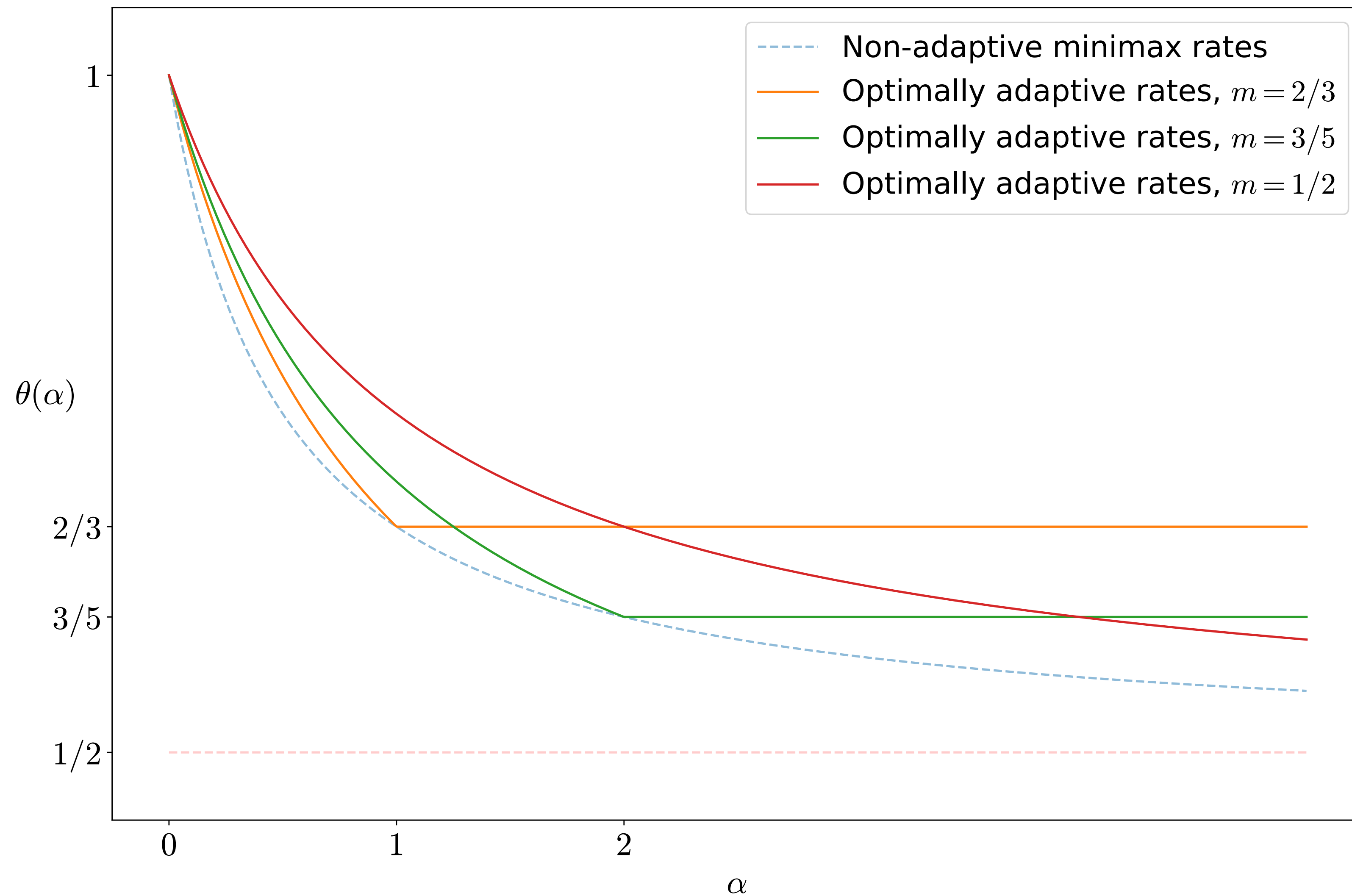
Answer: we cannot

## Theorem: Consequence of [Locatelli and Carpentier '18]

If a rate function  $\theta$  is achieved by some algorithm, then

$$\text{there exists } m \in [1/2, 1] \text{ s.t. } \theta(\alpha) \geq \max \left( m, 1 - m \frac{\alpha}{\alpha + 1} \right)$$

# Lower bound(s) on the adaptive rates



if  $R_T \leq c T^{\theta(\alpha)}$   
then  $\theta$  is lower bounded  
by one of these rates

Still, can we reach these rates?

# A bird's eye view of past approaches

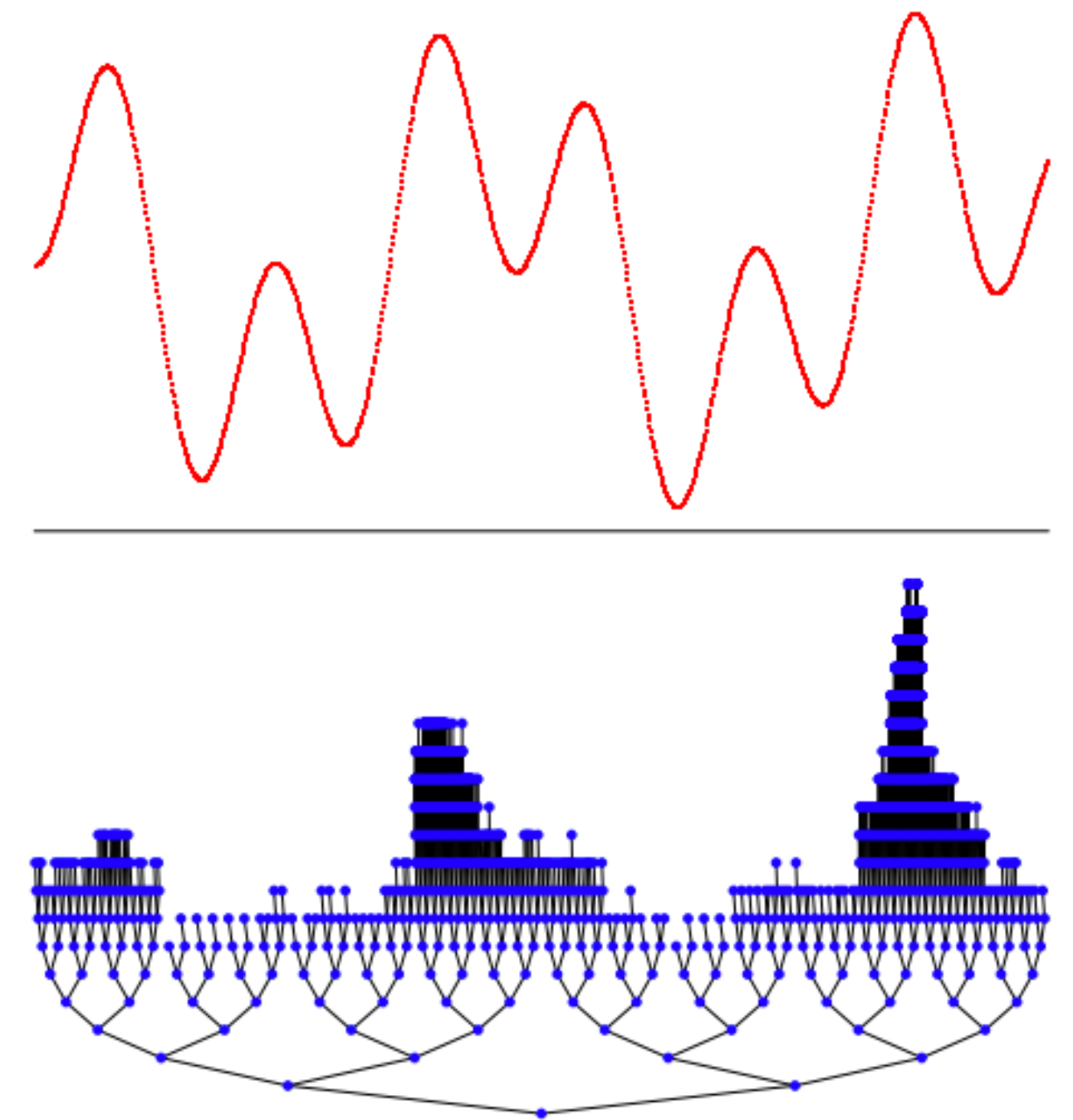
HOO [Bubeck, Munos, Stoltz, Szepesvári '11]

Zooming algorithm [Kleinberg, Upfal, Slivkins '13]

Adaptive-treed bandits [Bull '15]

SR [Locatelli, Carpentier '18]

e.g.



**Zoom in** on promising regions

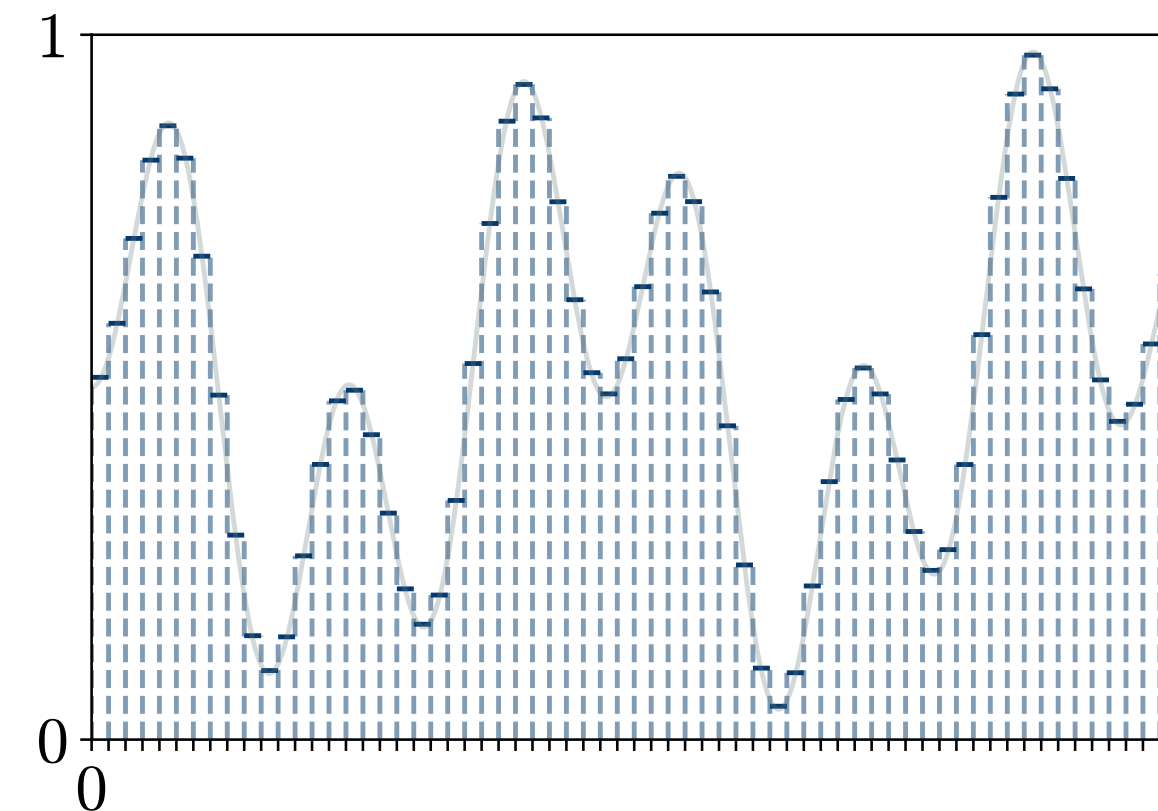
Select the promising regions **using the knowledge of the regularity**

# An optimally adaptive algorithm

## Basic idea

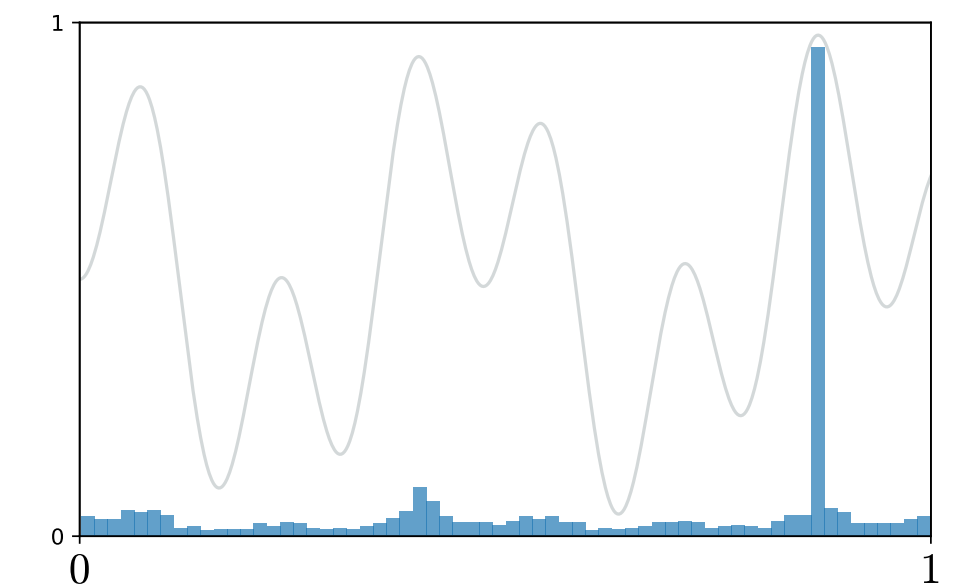
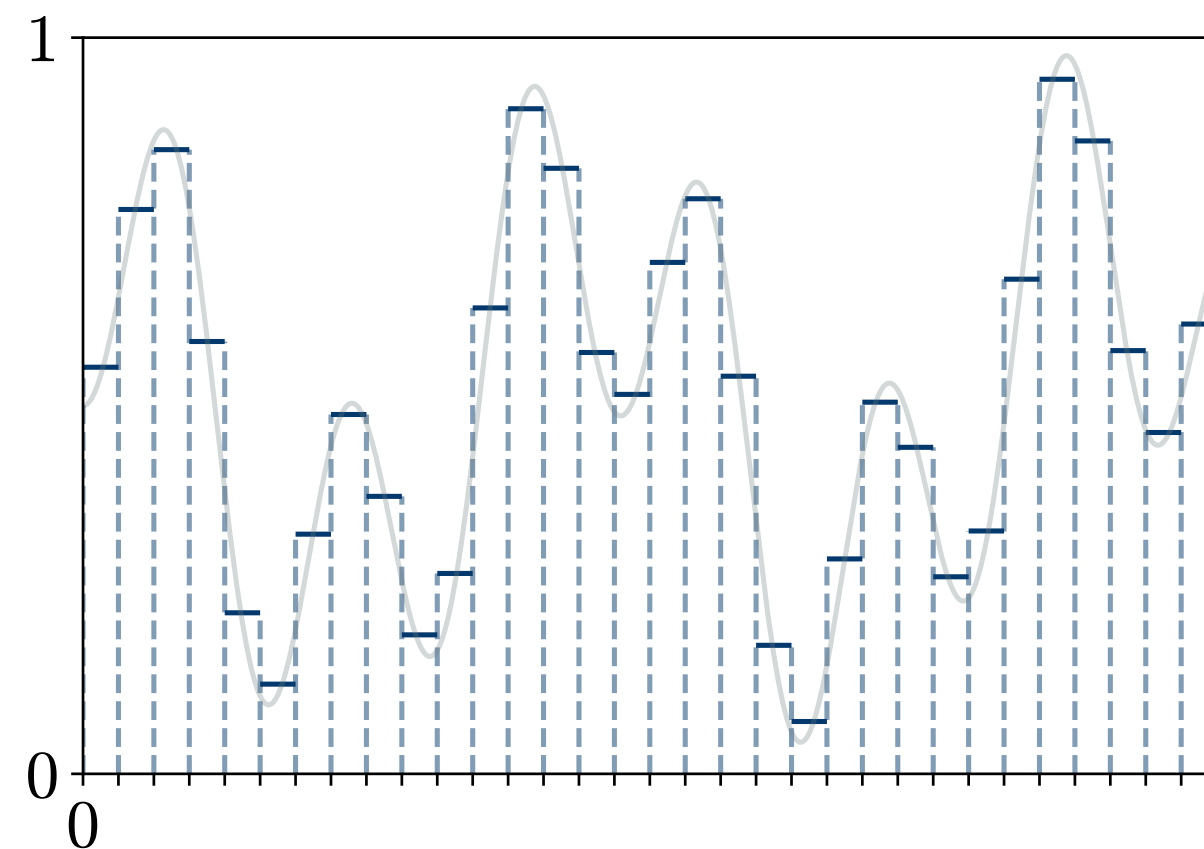
*Assume the worst...*

start with a very fine discretization  
but not for too long



*...then **ZOOM OUT.***

start over with a coarser discretization  
but remember what you played before



Discrete algorithm choosing between:

$K$  actions from the discretization

An extra-action: playing at random  
among actions selected in the past epoch

# Guarantees

## Algorithm : Memorize, Discretize, Zoom out

Set  $K_i \approx 2^{-i} \sqrt{T}$ ;  $D_i \approx 2^i \sqrt{T}$

For epochs 1 to  $\approx \log T$

– For  $D_i$  rounds, run a  $K_i$ -discretization **with memory of previous plays**

## Theorem: Medzo regret bound [H, 2019]

For any  $(L, \alpha)$ -Hölder function, without the knowledge of  $L$  or  $\alpha$ ,

$$R_T \leq \tilde{O} \left( L^{1/(\alpha+1)} T^{(\alpha+2)/(2\alpha+2)} \right)$$

with no assumption on  $\alpha$  and  $L$

# Conclusion

Adaptation leads to interesting questions with surprising phenomena

Achieving adaptation requires new algorithmic ideas

Ultimately important in practice/getting rid of hyperparameter tuning

**Bandits are fun!**

# Local perspectives

*KL-UCB*: Unifying notion of optimality in bandits?

*Range*:

- Minimax and optimal distribution-dependent rates for adapting to the lower end of the range
- Getting rid of the  $\sqrt{\log K}$

*Continuous bandits*:

- Including other types of regularity
- Is there a principled look at Medzo that could be applied elsewhere?
- What do we *need* to know to be able adapt at the usual rates [Locatelli Carpentier '18]?

# General perspectives

- What about non IID payoffs? [Zimmert and Seldin '19]
- Related problem 'Model selection in contextual multi-armed bandits'  
[Foster et al. '19], [Chatterji et al '19], [Foster et al. COLT open problem '20]

Observe context  $C_t \in \{1, \dots, S\}$  then choose  $A_t \in \{1, \dots, K\}$

A policy is mapping from context to action  $\pi : \mathcal{C} \rightarrow \mathcal{A}$

A model  $\Pi$  is a set of policies

Easy (Exp4):  $R_T(\text{Best } \pi \in \Pi) \leq c\sqrt{KT \log |\Pi|}$

Difficult: getting the same result with a sequence of nested models

-Dream general approach:

Can we combine a family of bandit algorithms and obtain one that is as good as the best?

[Agrawal et al. '17], [Pacchiano et al. '20], results are very specific

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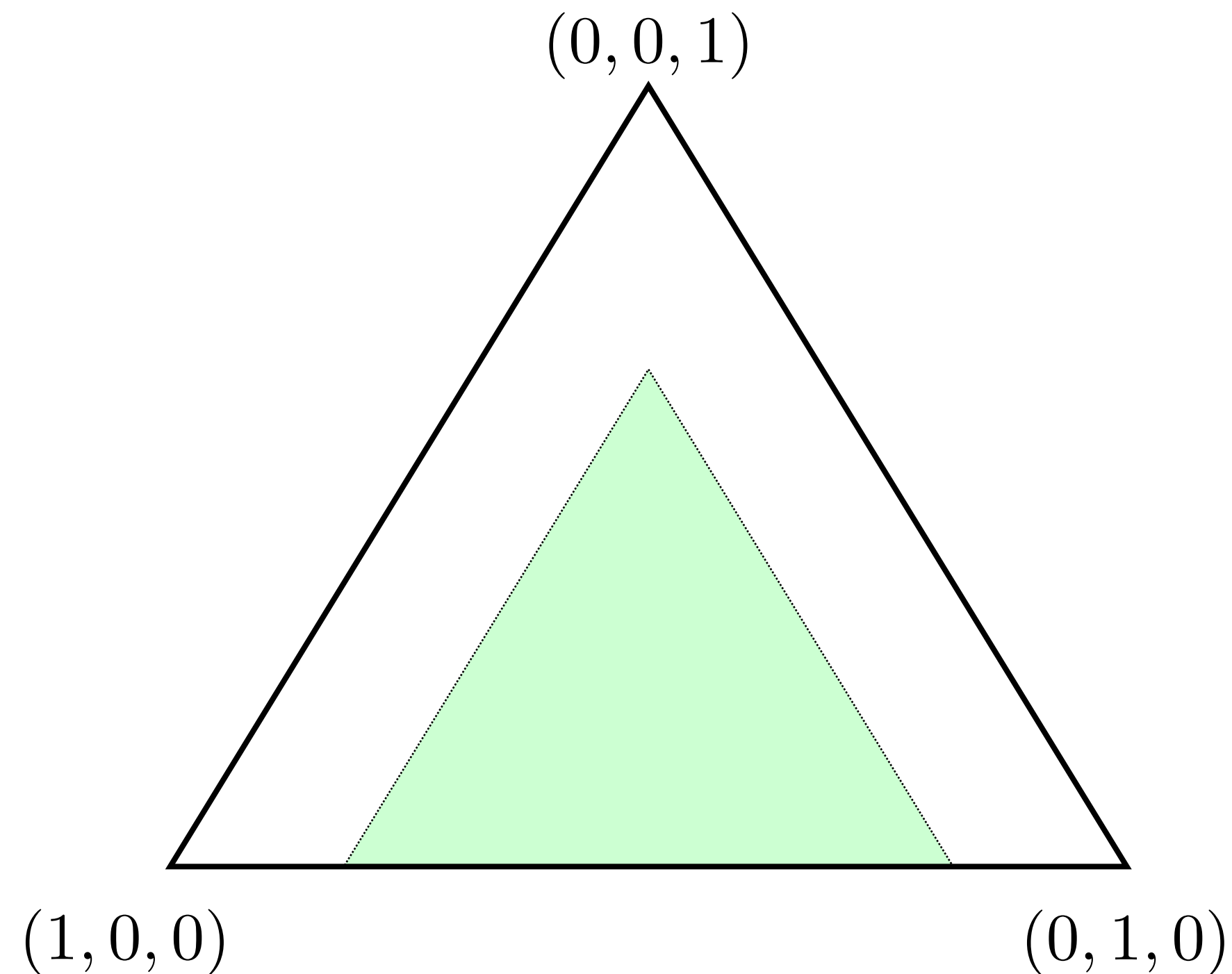
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# Bonus: diversity-preserving bandits

with Sébastien Gerchinovitz, Jean-Michel Loubes and Gilles Stoltz



Play a probability distribution  $p_t$  over  $\{1, 2, 3\}$

Observe  $A_t \sim p_t$ , and  $Y_t \sim \nu_{A_t}$

Require that  $p_{1,t}, p_{2,t} \geq 0.3$  (for example)

$$R_T = T \max_{p \text{ available}} \sum_{a=1}^K p_a \mu_a - \mathbb{E} \left[ \sum_{t=1}^T Y_t \right]$$

Bounded regret is possible  $\Leftrightarrow$  the best  $p$  is in the (relative) interior of the simplex

**Thank you Gilles and Pascal!**

**Thank you everyone!**